## 4 The Sperner property.

In this section we consider a surprising application of certain adjacency matrices to some problems in extremal set theory. An important role will also be played by finite groups. In general, extremal set theory is concerned with finding (or estimating) the most or least number of sets satisfying given settheoretic or combinatorial conditions. For example, a typical easy problem in extremal set theory is the following: What is the most number of subsets of an $n$-element set with the property that any two of them intersect? (Can you solve this problem?) The problems to be considered here are most conveniently formulated in terms of partially ordered sets, or posets for short. Thus we begin with discussing some basic notions concerning posets.
4.1 Definition. A poset (short for partially ordered set) $P$ is a finite set, also denoted $P$, together with a binary relation denoted $\leq$ satisfying the following axioms:
(P1) (reflexivity) $x \leq x$ for all $x \in P$
(P2) (antisymmetry) If $x \leq y$ and $y \leq x$, then $x=y$.
(P3) (transitivity) If $x \leq y$ and $y \leq z$, then $x \leq z$.

One easy way to obtain a poset is the following. Let $P$ be any collection of sets. If $x, y \in P$, then define $x \leq y$ in $P$ if $x \subseteq y$ as sets. It is easy to see that this definition of $\leq$ makes $P$ into a poset. If $P$ consists of all subsets of an $n$-element set $S$, then $P$ is called a (finite) boolean algebra of rank $n$ and is denoted by $B_{S}$. If $S=\{1,2, \ldots, n\}$, then we denote $B_{S}$ simply by $B_{n}$. Boolean algebras will play an important role throughout this section.

There is a simple way to represent small posets pictorially. The Hasse diagram of a poset $P$ is a planar drawing, with elements of $P$ drawn as dots. If $x<y$ in $P$ (i.e., $x \leq y$ and $x \neq y$ ), then $y$ is drawn "above" $x$ (i.e., with a larger vertical coordinate). An edge is drawn between $x$ and $y$ if $y$ covers $x$, i.e., $x<y$ and no element $z$ is in between, i.e., no $z$ satisfies $x<z<y$. By the transitivity property (P3), all the relations of a finite
poset are determined by the cover relations, so the Hasse diagram determines $P$. (This is not true for infinite posets; for instance, the real numbers $\mathbb{R}$ with their usual order is a poset with no cover relations.) The Hasse diagram of the boolean algebra $B_{3}$ looks like


We say that two posets $P$ and $Q$ are isomorphic if there is a bijection (one-to-one and onto function) $\varphi: P \rightarrow Q$ such that $x \leq y$ in $P$ if and only if $\varphi(x) \leq \varphi(y)$ in $Q$. Thus one can think that two posets are isomorphic if they differ only in the names of their elements. This is exactly analogous to the notion of isomorphism of groups, rings, etc. It is an instructive exercise to draw Hasse diagrams of the one poset of order (number of elements) one (up to isomorphism), the two posets of order two, the five posets of order three, and the sixteen posets of order four. More ambitious readers can try the 63 posets of order five, the 318 of order six, the 2045 of order seven, the 16999 of order eight, the 183231 of order nine, the 2567284 of order ten, the 46749427 of order eleven, the 1104891746 of order twelve, the 33823827452 of order thirteen, and the 1338193159771 of order fourteen. Beyond this the number is not currently known.

A chain $C$ in a poset is a totally ordered subset of $P$, i.e., if $x, y \in C$ then either $x \leq y$ or $y \leq x$ in $P$. A finite chain is said to have length $n$ if it has $n+1$ elements. Such a chain thus has the form $x_{0}<x_{1}<\cdots<x_{n}$. We say that a finite poset is graded of rank $n$ if every maximal chain has length $n$. (A chain is maximal if it's contained in no larger chain.) For instance, the boolean algebra $B_{n}$ is graded of rank $n$ [why?]. A chain $y_{0}<y_{1}<\cdots<y_{j}$ is said to be saturated if each $y_{i+1}$ covers $y_{i}$. Such a chain need not be maximal since there can be elements of $P$ smaller than $y_{0}$ or greater than $y_{j}$. If $P$ is graded of rank $n$ and $x \in P$, then we say that $x$ has rank $j$, denoted $\rho(x)=j$, if some (or equivalently, every) saturated chain of $P$ with top element $x$ has length $j$. Thus [why?] if we let $P_{j}=\{x \in P: \rho(x)=j\}$, then $P$ is a
disjoint union $P=P_{0} \cup P_{1} \cup \cdots \cup P_{n}$, and every maximal chain of $P$ has the form $x_{0}<x_{1}<\cdots<x_{n}$ where $\rho\left(x_{j}\right)=j$. We write $p_{j}=\left|P_{j}\right|$, the number of elements of $P$ of rank $j$. For example, if $P=B_{n}$ then $\rho(x)=|x|$ (the cardinality of $x$ as a set) and

$$
p_{j}=\#\{x \subseteq\{1,2, \cdots, n\}:|x|=j\}=\binom{n}{j} .
$$

(Note that we use both $|S|$ and $\# x$ for the cardinality of the finite set $S$.)

We say that a graded poset $P$ of rank $n$ (always assumed to be finite) is rank-symmetric if $p_{i}=p_{n-i}$ for $0 \leq i \leq n$, and rank-unimodal if $p_{0} \leq$ $p_{1} \leq \cdots \leq p_{j} \geq p_{j+1} \geq p_{j+2} \geq \cdots \geq p_{n}$ for some $0 \leq j \leq n$. If $P$ is both rank-symmetric and rank-unimodal, then we clearly have

$$
\begin{gathered}
p_{0} \leq p_{1} \leq \cdots \leq p_{m} \geq p_{m+1} \geq \cdots \geq p_{n}, \text { if } n=2 m \\
p_{0} \leq p_{1} \leq \cdots \leq p_{m}=p_{m+1} \geq p_{m+2} \geq \cdots \geq p_{n}, \text { if } n=2 m+1 .
\end{gathered}
$$

We also say that the sequence $p_{0}, p_{1}, \ldots, p_{n}$ itself or the polynomial $F(q)=$ $p_{0}+p_{1} q+\cdots+p_{n} q^{n}$ is symmetric or unimodal, as the case may be. For instance, $B_{n}$ is rank-symmetric and rank-unimodal, since it is well-known (and easy to prove) that the sequence $\binom{n}{0},\binom{n}{1}, \ldots,\binom{n}{n}$ (the $n$th row of Pascal's triangle) is symmetric and unimodal. Thus the polynomial $(1+q)^{n}$ is symmetric and unimodal.

A few more definitions, and then finally some results! An antichain in a poset $P$ is a subset $A$ of $P$ for which no two elements are comparable, i.e., we can never have $x, y \in A$ and $x<y$. For instance, in a graded poset $P$ the "levels" $P_{j}$ are antichains [why?]. We will be concerned with the problem of finding the largest antichain in a poset. Consider for instance the boolean algebra $B_{n}$. The problem of finding the largest antichain in $B_{n}$ is clearly equivalent to the following problem in extremal set theory: Find the largest collection of subsets of an $n$-element set such that no element of the collection contains another. A good guess would be to take all the subsets of cardinality $\lfloor n / 2\rfloor$ (where $\lfloor x\rfloor$ denotes the greatest integer $\leq x$ ), giving a total of $\binom{n}{\lfloor n / 2\rfloor}$ sets in all. But how can we actually prove there is no larger collection? Such a proof was first given by Emmanuel Sperner in 1927 and is known as Sperner's theorem. We will give two proofs of Sperner's theorem
in this section; one proof uses linear algebra and will be applied to certain other situations, while the other proof is an elegant combinatorial argument due to David Lubell in 1966, which we present for its "cultural value." Our extension of Sperner's theorem to certain other situations will involve the following crucial definition.
4.2 Definition. Let $P$ be a graded poset of rank $n$. We say that $P$ has the Sperner property or is a Sperner poset if

$$
\max \{|A|: A \text { is an antichain of } P\}=\max \left\{\left|P_{i}\right|: 0 \leq i \leq n\right\}
$$

In other words, no antichain is larger than the largest level $P_{i}$.
Thus Sperner's theorem is equivalent to saying that $B_{n}$ has the Sperner property. Note that if $P$ has the Sperner property there may still be antichains of maximum cardinality other than the biggest $P_{i}$; there just can't be any bigger antichains.
4.3 Example. A simple example of a graded poset that fails to satisfy the Sperner property is the following:


We now will discuss a simple combinatorial condition which guarantees that certain graded posets $P$ are Sperner. We define an order-matching from $P_{i}$ to $P_{i+1}$ to be a one-to-one function $\mu: P_{i} \rightarrow P_{i+1}$ satisfying $x<\mu(x)$ for all $x \in P_{i}$. Clearly if such an order-matching exists then $p_{i} \leq p_{i+1}$ (since $\mu$ is one-to-one). Easy examples show that the converse is false, i.e., if $p_{i} \leq p_{i+1}$ then there need not exist an order-matching from $P_{i}$ to $P_{i+1}$. We similarly define an order-matching from $P_{i}$ to $P_{i-1}$ to be a one-to-one function $\mu: P_{i} \rightarrow P_{i-1}$ satisfying $\mu(x)<x$ for all $x \in P_{i}$.
4.4 Proposition. Let $P$ be a graded poset of rank n. Suppose there exists an integer $0 \leq j \leq n$ and order-matchings

$$
\begin{equation*}
P_{0} \rightarrow P_{1} \rightarrow P_{2} \rightarrow \cdots \rightarrow P_{j} \leftarrow P_{j+1} \leftarrow P_{j+2} \leftarrow \cdots \leftarrow P_{n} . \tag{17}
\end{equation*}
$$

Then $P$ is rank-unimodal and Sperner.

Proof. Since order-matchings are one-to-one it is clear that

$$
p_{0} \leq p_{1} \leq \cdots \leq p_{j} \geq p_{j+1} \geq p_{j+2} \geq \cdots \geq p_{n}
$$

Hence $P$ is rank-unimodal.
Define a graph $G$ as follows. The vertices of $G$ are the elements of $P$. Two vertices $x, y$ are connected by an edge if one of the order-matchings $\mu$ in the statement of the proposition satisfies $\mu(x)=y$. (Thus $G$ is a subgraph of the Hasse diagram of $P$.) Drawing a picture will convince you that $G$ consists of a disjoint union of paths, including single-vertex paths not involved in any of the order-matchings. The vertices of each of these paths form a chain in $P$. Thus we have partitioned the elements of $P$ into disjoint chains. Since $P$ is rank-unimodal with biggest level $P_{j}$, all of these chains must pass through $P_{j}$ [why?]. Thus the number of chains is exactly $p_{j}$. Any antichain $A$ can intersect each of these chains at most once, so the cardinality $|A|$ of $A$ cannot exceed the number of chains, i.e., $|A| \leq p_{j}$. Hence by definition $P$ is Sperner.

It is now finally time to bring some linear algebra into the picture. For any (finite) set $S$, we let $\mathbb{R} S$ denote the real vector space consisting of all formal linear combinations (with real coefficients) of elements of $S$. Thus $S$ is a basis for $\mathbb{R} S$, and in fact we could have simply defined $\mathbb{R} S$ to be the real vector space with basis $S$. The next lemma relates the combinatorics we have just discussed to linear algebra and will allow us to prove that certain posets are Sperner by the use of linear algebra (combined with some finite group theory).
4.5 Lemma. Suppose there exists a linear transformation $U: \mathbb{R} P_{i} \rightarrow$ $\mathbb{R} P_{i+1}$ ( $U$ stands for "up") satisfying:

- $U$ is one-to-one.
- For all $x \in P_{i}, U(x)$ is a linear combination of elements $y \in P_{i+1}$ satisfying $x<y$. (We then call $U$ an order-raising operator.)

Then there exists an order-matching $\mu: P_{i} \rightarrow P_{i+1}$.

Similarly, suppose there exists a linear transformation $U: \mathbb{R} P_{i} \rightarrow \mathbb{R} P_{i+1}$ satisfying:

- $U$ is onto.
- $U$ is an order-raising operator.

Then there exists an order-matching $\mu: P_{i+1} \rightarrow P_{i}$.
Proof. Suppose $U: \mathbb{R} P_{i} \rightarrow \mathbb{R} P_{i+1}$ is a one-to-one order-raising operator. Let $[U]$ denote the matrix of $U$ with respect to the bases $P_{i}$ of $\mathbb{R} P_{i}$ and $P_{i+1}$ of $\mathbb{R} P_{i+1}$. Thus the columns of $[U]$ are indexed by the elements $x_{1}, \ldots, x_{p_{i}}$ of $P_{i}$ (in some order) and the rows by the elements $y_{1}, \ldots, y_{p_{i+1}}$ of $P_{i+1}$. Since $U$ is one-to-one, the rank of $[U]$ is equal to $p_{i}$ (the number of columns). Since the column rank of a matrix equals its row rank, $[U]$ must have $p_{i}$ linearly independent rows. Say we have labelled the elements of $P_{i+1}$ so that the first $p_{i}$ rows of $[U]$ are linearly independent.

Let $A=\left(a_{i j}\right)$ be the $p_{i} \times p_{i}$ matrix whose rows are the first $p_{i}$ rows of $[U]$. (Thus $A$ is a square submatrix of $[U]$.) Since the rows of $A$ are linearly independent, we have

$$
\operatorname{det}(A)=\sum \pm a_{\pi(1), 1} \cdots a_{\pi\left(p_{i}\right), p_{i}} \neq 0
$$

where the sum is over all permutations $\pi$ of $1, \ldots, p_{i}$. Thus some term $\pm a_{\pi(1), 1} \cdots a_{\pi\left(p_{i}\right), p_{i}}$ of the above sum in nonzero. Since $U$ is order-raising, this means that [why?] $x_{k}<y_{\pi(k)}$ for $1 \leq k \leq p_{i}$. Hence the map $\mu: P_{i} \rightarrow P_{i+1}$ defined by $\mu\left(x_{k}\right)=y_{\pi(k)}$ is an order-matching, as desired.

The case when $U$ is onto rather than one-to-one is proved by a completely analogous argument.

We now want to apply Proposition 4.4 and Lemma 4.5 to the boolean algebra $B_{n}$. For each $0 \leq i<n$, we need to define a linear transformation $U_{i}: \mathbb{R}\left(B_{n}\right)_{i} \rightarrow \mathbb{R}\left(B_{n}\right)_{i+1}$, and then prove it has the desired properties. We simply define $U_{i}$ to be the simplest possible order-raising operator, namely,
for $x \in\left(B_{n}\right)_{i}$, let

$$
\begin{equation*}
U_{i}(x)=\sum_{\substack{y \in(B n)_{i+1} \\ y>x}} y . \tag{18}
\end{equation*}
$$

Note that since $\left(B_{n}\right)_{i}$ is a basis for $\mathbb{R}\left(B_{n}\right)_{i}$, equation (18) does indeed define a unique linear transformation $U_{i}: \mathbb{R}\left(B_{n}\right)_{i} \rightarrow \mathbb{R}\left(B_{n}\right)_{i+1}$. By definition $U_{i}$ is order-raising; we want to show that $U_{i}$ is one-to-one for $i<n / 2$ and onto for $i \geq n / 2$. There are several ways to show this using only elementary linear algebra; we will give what is perhaps the simplest proof, though it is quite tricky. The idea is to introduce "dual" operators $D_{i}: \mathbb{R}\left(B_{n}\right)_{i} \rightarrow\left(B_{n}\right)_{i-1}$ to the $U_{i}$ 's ( $D$ stands for "down"), defined by

$$
\begin{equation*}
D_{i}(y)=\sum_{\substack{x \in\left(B_{n}\right)_{i-1} \\ x<y}} x \tag{19}
\end{equation*}
$$

for all $y \in\left(B_{n}\right)_{i}$. Let $\left[U_{i}\right]$ denote the matrix of $U_{i}$ with respect to the bases $\left(B_{n}\right)_{i}$ and $\left(B_{n}\right)_{i+1}$, and similarly let $\left[D_{i}\right]$ denote the matrix of $D_{i}$ with respect to the bases $\left(B_{n}\right)_{i}$ and $\left(B_{n}\right)_{i-1}$. A key observation which we will use later is that

$$
\begin{equation*}
\left[D_{i+1}\right]=\left[U_{i}\right]^{t} \tag{20}
\end{equation*}
$$

i.e., the matrix $\left[D_{i+1}\right.$ ] is the transpose of the matrix $\left[U_{i}\right.$ ] [why?]. Now let $I_{i}: \mathbb{R}\left(B_{n}\right)_{i} \rightarrow \mathbb{R}\left(B_{n}\right)_{i}$ denote the identity transformation on $\mathbb{R}\left(B_{n}\right)_{i}$, i.e., $I_{i}(u)=u$ for all $u \in \mathbb{R}\left(B_{n}\right)_{i}$. The next lemma states (in linear algebraic terms) the fundamental combinatorial property of $B_{n}$ which we need. For this lemma set $U_{n}=0$ and $D_{0}=0$ (the 0 linear transformation between the appropriate vector spaces).
4.6 Lemma. Let $0 \leq i \leq n$. Then

$$
\begin{equation*}
D_{i+1} U_{i}-U_{i-1} D_{i}=(n-2 i) I_{i} . \tag{21}
\end{equation*}
$$

(Linear transformations are multiplied right-to-left, so $A B(u)=A(B(u))$.)
Proof. Let $x \in\left(B_{n}\right)_{i}$. We need to show that if we apply the left-hand side of (21) to $x$, then we obtain $(n-2 i) x$. We have

$$
D_{i+1} U_{i}(x)=D_{i+1}\left(\sum_{\substack{|y|=i+1 \\ x \subset y}} y\right)
$$

$$
=\sum_{\substack{|y|=i+1 \\ x \subset y}} \sum_{\substack{|z|=i \\ z \subset y}} z
$$

If $x, z \in\left(B_{n}\right)_{i}$ satisfy $|x \cap z|<i-1$, then there is no $y \in\left(B_{n}\right)_{i+1}$ such that $x \subset y$ and $z \subset y$. Hence the coefficient of $z$ in $D_{i+1} U_{i}(x)$ when it is expanded in terms of the basis $\left(B_{n}\right)_{i}$ is 0 . If $|x \cap z|=i-1$, then there is one such $y$, namely, $y=x \cup z$. Finally if $x=z$ then $y$ can be any element of $\left(B_{n}\right)_{i+1}$ containing $x$, and there are $n-i$ such $y$ in all. It follows that

$$
\begin{equation*}
D_{i+1} U_{i}(x)=(n-i) x+\sum_{\substack{|z z=i\\| x \cap \mid=i-1}} z . \tag{22}
\end{equation*}
$$

By exactly analogous reasoning (which the reader should check), we have for $x \in\left(B_{n}\right)_{i}$ that

$$
\begin{equation*}
U_{i-1} D_{i}(x)=i x+\sum_{\substack{|z|=i \\|x \cap z|=i-1}} z . \tag{23}
\end{equation*}
$$

Subtracting (23) from (22) yields $\left(D_{i+1} U_{i}-U_{i-1} D_{i}\right)(x)=(n-2 i) x$, as desired.
4.7 Theorem. The operator $U_{i}$ defined above is one-to-one if $i<n / 2$ and is onto if $i \geq n / 2$.

Proof. Recall that $\left[D_{i}\right]=\left[U_{i-1}\right]^{t}$. From linear algebra we know that a (rectangular) matrix times its transpose is positive semidefinite (or just semidefinite for short) and hence has nonnegative (real) eigenvalues. By Lemma 4.6 we have

$$
D_{i+1} U_{i}=U_{i-1} D_{i}+(n-2 i) I_{i} .
$$

Thus the eigenvalues of $D_{i+1} U_{i}$ are obtained from the eigenvalues of $U_{i-1} D_{i}$ by adding $n-2 i$. Since we are assuming that $n-2 i>0$, it follows that the eigenvalues of $D_{i+1} U_{i}$ are strictly positive. Hence $D_{i+1} U_{i}$ is invertible (since it has no 0 eigenvalues). But this implies that $U_{i}$ is one-to-one [why?], as desired.

The case $i \geq n / 2$ is done by a "dual" argument (or in fact can be deduced directly from the $i<n / 2$ case by using the fact that the poset $B_{n}$ is "selfdual," though we will not go into this). Namely, from the fact that

$$
U_{i} D_{i+1}=D_{i+2} U_{i+1}+(2 i+2-n) I_{i+1}
$$

we get that $U_{i} D_{i+1}$ is invertible, so now $U_{i}$ is onto, completing the proof.
Combining Proposition 4.4, Lemma 4.5, and Theorem 4.7, we obtain the famous theorem of Sperner.
4.8 Corollary. The boolean algebra $B_{n}$ has the Sperner property.

It is natural to ask whether there is a less indirect proof of Corollary 4.8. In fact, several nice proofs are known; we give one due to David Lubell, mentioned before Definition 4.2.

Lubell's proof of Sperner's theorem. First we count the total number of maximal chains $\emptyset=x_{0}<x_{1}<\cdots<x_{n}=\{1, \ldots, n\}$ in $B_{n}$. There are $n$ choices for $x_{1}$, then $n-1$ choices for $x_{2}$, etc., so there are $n$ ! maximal chains in all. Next we count the number of maximal chains $x_{0}<x_{1}<\cdots<x_{i}=$ $x<\cdots<x_{n}$ which contain a given element $x$ of rank $i$. There are $i$ choices for $x_{1}$, then $i-1$ choices for $x_{2}$, up to one choice for $x_{i}$. Similarly there are $n-i$ choices for $x_{i+1}$, then $n-2$ choices for $x_{i+2}$, etc., up to one choice for $x_{n}$. Hence the number of maximal chains containing $x$ is $i!(n-i)$ !.

Now let $A$ be an antichain. If $x \in A$, then let $C_{x}$ be the set of maximal chains of $B_{n}$ which contain $x$. Since $A$ is an antichain, the sets $C_{x}, x \in A$ are pairwise disjoint. Hence

$$
\begin{aligned}
\left|\bigcup_{x \in A} C_{x}\right| & =\sum_{x \in A}\left|C_{x}\right| \\
& =\sum_{x \in A}(\rho(x))!(n-\rho(x))!
\end{aligned}
$$

Since the total number of maximal chains in the $C_{x}$ 's cannot exceed the total number $n$ ! of maximal chains in $B_{n}$, we have

$$
\sum_{x \in A}(\rho(x))!(n-\rho(x))!\leq n!
$$

Divide both sides by $n$ ! to obtain

$$
\sum_{x \in A} \frac{1}{\binom{n}{\rho(x)}} \leq 1
$$

Since $\binom{n}{i}$ is maximized when $i=\lfloor n / 2\rfloor$, we have

$$
\frac{1}{\binom{n}{\lfloor n / 2\rfloor}} \leq \frac{1}{\binom{n}{\rho(x)}}
$$

for all $x \in A$ (or all $x \in B_{n}$ ). Thus

$$
\sum_{x \in A} \frac{1}{\binom{n}{\lfloor n / 2\rfloor}} \leq 1
$$

or equivalently,

$$
|A| \leq\binom{ n}{\lfloor n / 2\rfloor}
$$

Since $\binom{n}{\lfloor n / 2\rfloor}$ is the size of the largest level of $B_{n}$, it follows that $B_{n}$ is Sperner.

In view of the above elegant proof of Lubell, the reader may be wondering what was the point of giving a rather complicated and indirect proof using linear algebra. Admittedly, if all we could obtain from the linear algebra machinery we have developed was just another proof of Sperner's theorem, then it would have been hardly worth the effort. But in the next section we will show how Theorem 4.7, when combined with a little finite group theory, can be used to obtain many interesting combinatorial results for which simple, direct proofs are not known.

